

Advanced Robot Control

Real-Time Operating System

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Presentation compiled for taking notes during lecture



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- 1 Operating system
 - OS support features
 - Real-Time Operating System
 - Shadowed stack pointer
 - Supervised and unsupervised mode
 - Context switching
 - Exclusive accesses
- 2 FreeRTOS
 - Introduction
 - Tasks
 - Mutexes
 - Queues
 - Event groups
 - Software timers
 - Heap management



Operating System (1/2)

What is an operating system?



Operating System (2/2)

It is a specific software that runs in kernel mode (supervised mode) [3].



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- 3 **Quiz**



OS support features in Cortex-M3/M4 core (1/1)

Cortex-M3/M4 support following features to facilitate the implementation of an embedded OS [4]:

- **Shadowed stack pointer**
- **SysTick timer**
- **Supervisor Call (SVC) and Pendable Service Call (PendSV) exceptions**
- **Unprivileged execution level**
- **Exclusive accesses**



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Real-Time Operating System

The RTOS is a special kind of an operating system which focuses on time. Two groups of RTOS can be specified:

- soft real-time operating system,
- hard real-time operating system.



Hard RTOS

Hard real-time operating system is a specific branch of RTOS which has to meet strict time constraints.



Soft RTOS

In contrast to Hard RTOS a Soft RTOS is an operating system where occasional violation of strict time regime is undesirable but will not lead to any permanent damage.



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Main Stack Pointer and Processor Stack Pointer

Figure: Stack pointer in a task [4]



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Memory Protection Unit (1/2)

The MPU is a programmable device that can be used to define memory access permissions (e.g., privileged access only or full access) and memory attributes (e.g., bufferable, cacheable) for different memory regions [4].



Memory Protection Unit (2/2)

The MPU can be used to make an embedded system more robust .



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Starting task with SVC

Figure: Initialization of a task [4]



Switching task with PendSV

Figure: Context switching [4]



Switching task during interrupt (1/2)

Figure: Improper task switching [4]



Switching task during interrupt (2/2)

Figure: Proper task switching [4]



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Hardware Mutex support (1/1)

A common way to limit the access to a resource is a semaphore. It can be represented as a counter which holds a number of resource.



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Introduction

FreeRTOS is a real-time kernel on top of which embedded applications can be built to meet their hard real-time requirements [1].



Features (1/1)

- **Abstracting away timing information**
- **Maintainability/Extensibility**
- **Modularity**
- **Team development**
- **Easier testing**
- **Code reuse**
- **Improved efficiency**
- **Idle time**
- **Power Management**
- **Flexible interrupt handling**
- **Mixed processing requirements**



ISR dedicated API

ISR specific API

It is worthy of note that not all API can be called in ISR routines.



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Task states

Figure: Task switching [1]



Task creation

```
1 BaseType_t xTaskCreate( TaskFunction_t pvTaskCode ,
2     const char * const pcName,
3     uint16_t usStackDepth ,
4     void *pvParameters ,
5     UBaseType_t uxPriority ,
6     TaskHandle_t *pxCreatedTask );

1 vTaskStartScheduler ();
```



Delaying task execution

```
1 void vTaskDelay( TickType_t xTicksToDelay );

1 void vTaskDelayUntil( TickType_t * pxPreviousWakeTime,
2                       TickType_t xTimeIncrement );

1 TickType_t xLastWakeTime;
2 xLastWakeTime = xTaskGetTickCount();
```



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Mutexes (1/2)

A mutex can be created using [2]

- 1 SemaphoreHandle_t xSemaphoreCreateMutex(void);
- 1 BaseType_t xSemaphoreTake(SemaphoreHandle_t xSemaphore,
TickType_t xTicksToWait);
- 1 BaseType_t xSemaphoreGive(SemaphoreHandle_t xSemaphore);



Mutexes (2/2)

During usage of mutexes a situation called *Deadlock* can arise.



Critical section

A critical section is a region of code surrounded with

```
1 taskENTER_CRITICAL ();  
2 //...  
3 taskEXIT_CRITICAL ();
```



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Queues (1/3)

Queues are normally used as **First In First Out (FIFO) buffers**, where data is written to the end (tail) of the queue and removed from the front (head) of the queue.

```
1 QueueHandle_t xQueueCreate( UBaseType_t uxQueueLength,  
2 UBaseType_t uxItemSize );
```

```
1 BaseType_t xQueueSendToFront( QueueHandle_t xQueue,  
2 const void * pvItemToQueue,  
3 TickType_t xTicksToWait );
```

```
1 BaseType_t xQueueSendToBack( QueueHandle_t xQueue,  
2 const void * pvItemToQueue,  
3 TickType_t xTicksToWait );
```



Queues (2/3)

```
1 BaseType_t xQueueReceive( QueueHandle_t xQueue,  
2 void * const pvBuffer,  
3 TickType_t xTicksToWait );  
  
1 UBaseType_t uxQueueMessagesWaiting( QueueHandle_t xQueue );
```



Queues (3/3)

Queue can be used as a **message box** where size of a queue is equal to 1.

```
1 BaseType_t xQueueOverwrite( QueueHandle_t xQueue,  
2 const void * pvItemToQueue );
```

```
1 BaseType_t xQueuePeek( QueueHandle_t xQueue,  
2 void * const pvBuffer,  
3 TickType_t xTicksToWait );
```



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Event groups (1/2)

Event group is a mechanism available in FreeRTOS which allows to suspend a task in a blocking state as long as a certain condition is not met.



Event groups (2/2)

```
1 EventGroupHandle_t xEventGroupCreate( void );

1 EventBits_t xEventGroupSetBits( EventGroupHandle_t xEventGroup ,
2 const EventBits_t uxBitsToSet );

1 EventBits_t xEventGroupWaitBits( const EventGroupHandle_t
    xEventGroup ,
2 const EventBits_t uxBitsToWaitFor ,
3 const BaseType_t xClearOnExit ,
4 const BaseType_t xWaitForAllBits ,
5 TickType_t xTicksToWait );
```



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Software timers (1/4)

Software timers can run a function (callback function) at a specific time or periodically.



Software timers (2/4)

```
1 TimerHandle_t xTimerCreate( const char * const pcTimerName,  
2 TickType_t xTimerPeriodInTicks,  
3 UBaseType_t uxAutoReload,  
4 void * pvTimerID,  
5 TimerCallbackFunction_t pxCallbackFunction );  
  
1 BaseType_t xTimerStart( TimerHandle_t xTimer, TickType_t  
    xTicksToWait );  
  
1 BaseType_t xTimerStop( TimerHandle_t xTimer, TickType_t  
    xTicksToWait );
```



Software timers (3/4)

```
1 void vTimerSetTimerID( const TimerHandle_t xTimer, void *pvNewID
    );

1 void *pvTimerGetTimerID( TimerHandle_t xTimer );
```



Software timers (4/4)

```
1 BaseType_t xTimerChangePeriod( TimerHandle_t xTimer ,
2 TickType_t xNewTimerPeriodInTicks ,
3 TickType_t xTicksToWait );

1 BaseType_t xTimerReset( TimerHandle_t xTimer , TickType_t
    xTicksToWait );
```



State machine for one shoot sw timer (1/1)

Figure: One-shot software timer states and transitions [1]



State machine for auto-reload sw timer (1/1)

Figure: Auto-reload software timer states and transitions [1]



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Heap implementations

In FreeRTOS a number of heap implementations is available:

- **Heap_1**
- **Heap_2**
- **Heap_3**
- **Heap_4**
- **Heap_5**



Quiz (1/1)

Calculate group number as the rest from dividing the Student ID number by 4.

Example

Student ID number is 123456, thus the group is 0.

Take last 2 digits from Student ID number (56) and calculate the rest from dividing by 4 ($56 \% 4 = 0$).

Write down your name, Student ID number and group.



Literature (1/2)



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Literature (2/2)



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